

Deductive Verification of Data Structures

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Introduction

- I will mostly talk about a particular and well known data type

stack

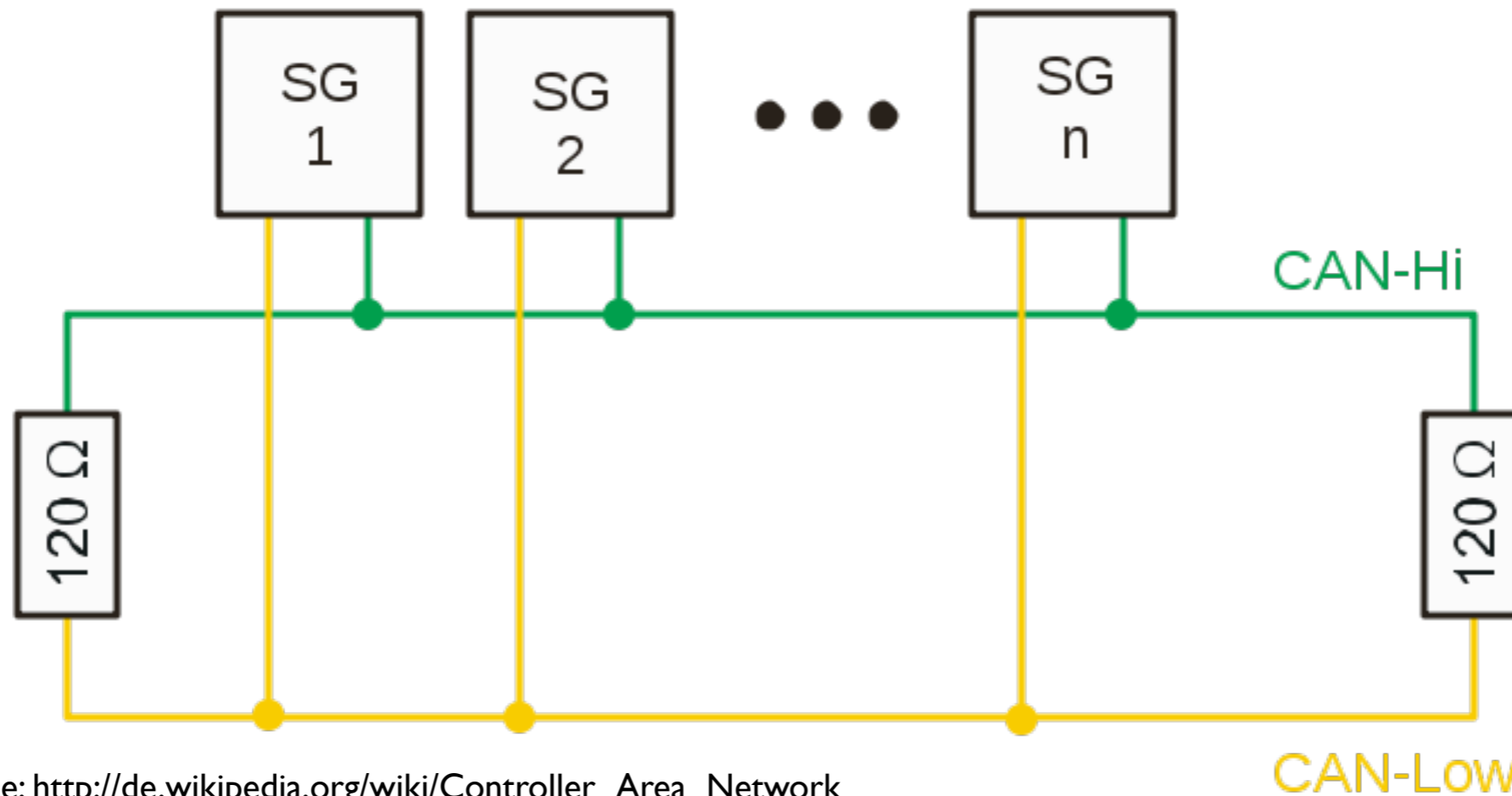
- The principles to be discussed can be applied to other data structures as well

Overview

- Motivation
- A simple C stack
- From unit tests to unit proofs
- Representation independence
- Conclusions

Motivation

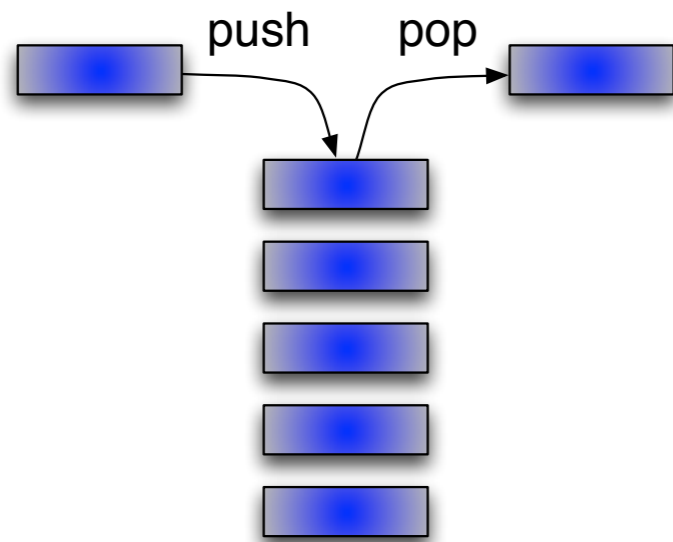
Vehicle Networks



Source: http://de.wikipedia.org/wiki/Controller_Area_Network

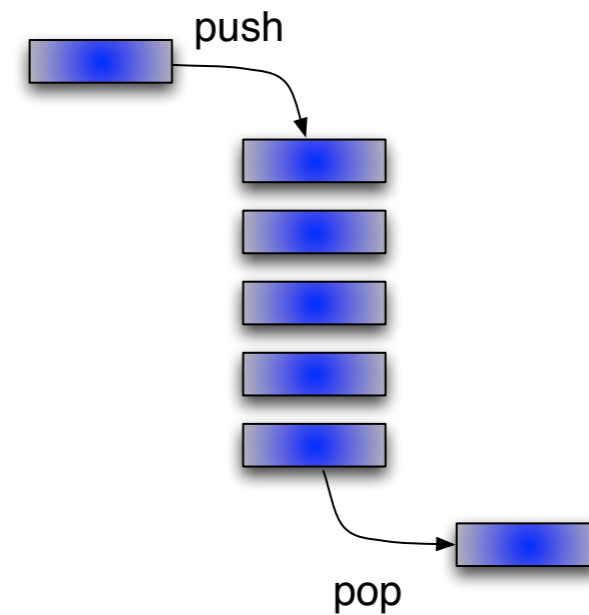
Buffering of Messages

Stack



first in, last out

Queue



first in, first out

C++ standard library

- `std::stack<T>` and `std::queue<T>`
 - generic and type safe
 - portable and operationally proven
- disadvantages (for embedded systems)
 - not C, dynamic memory

A simple C Stack

Design guide lines

- interface similar to `std::stack`
- can only hold a maximum number of elements
- no need for dynamic memory

What exactly is a stack?

- a data type with associated functions
 - `create`
 - `empty, full`
 - `push, top, pop,`
- functions must satisfy *stack axioms*

Some stack axioms

`IsEmpty(create()) = TRUE`

`IsEmpty(push(x,s)) = FALSE`

`pop(push(x,s)) = s`

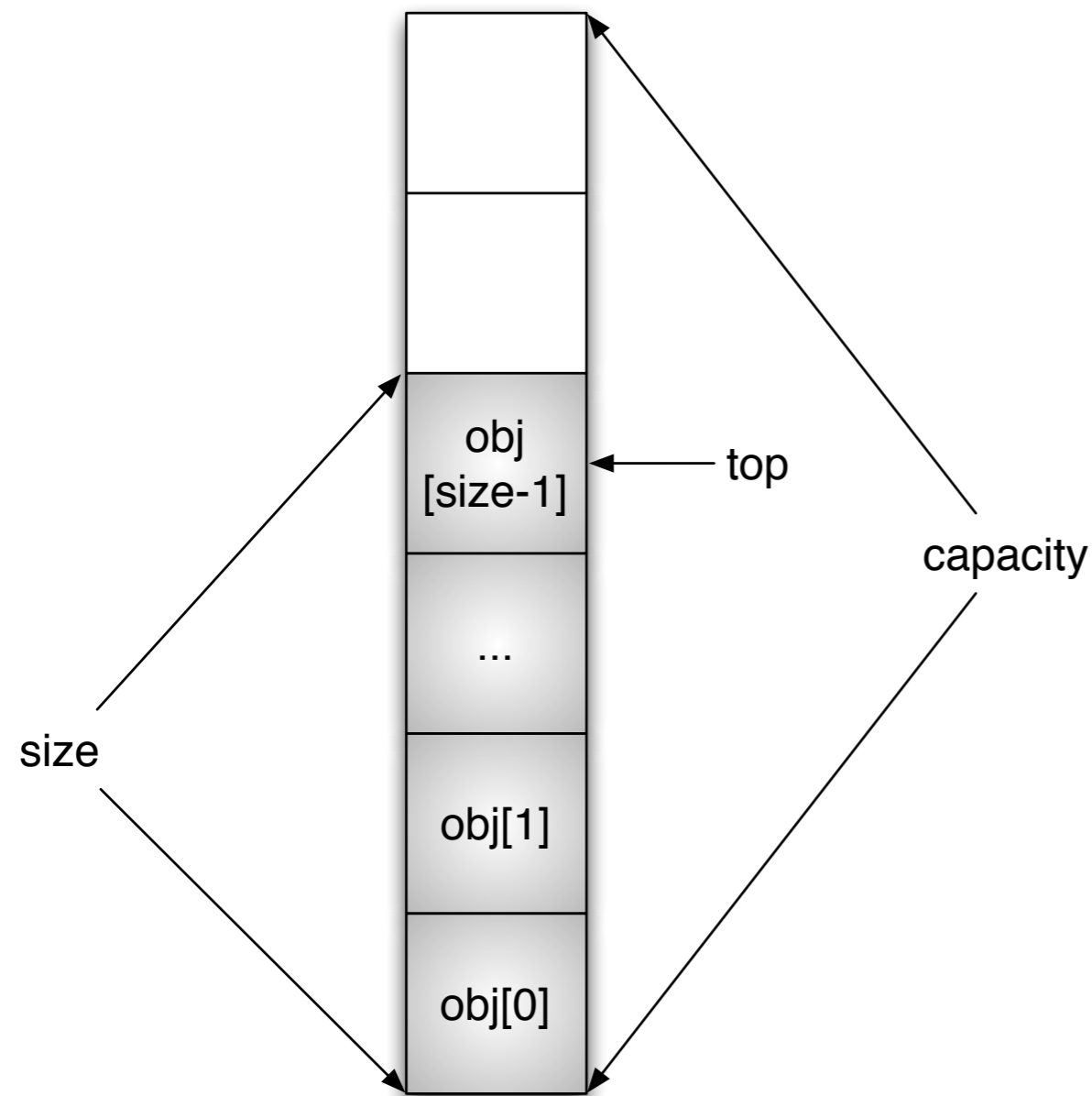
`top(push(x,s)) = x`

`push(top(s),pop(s)) = s, if s not empty`

Focus on Axiom 3

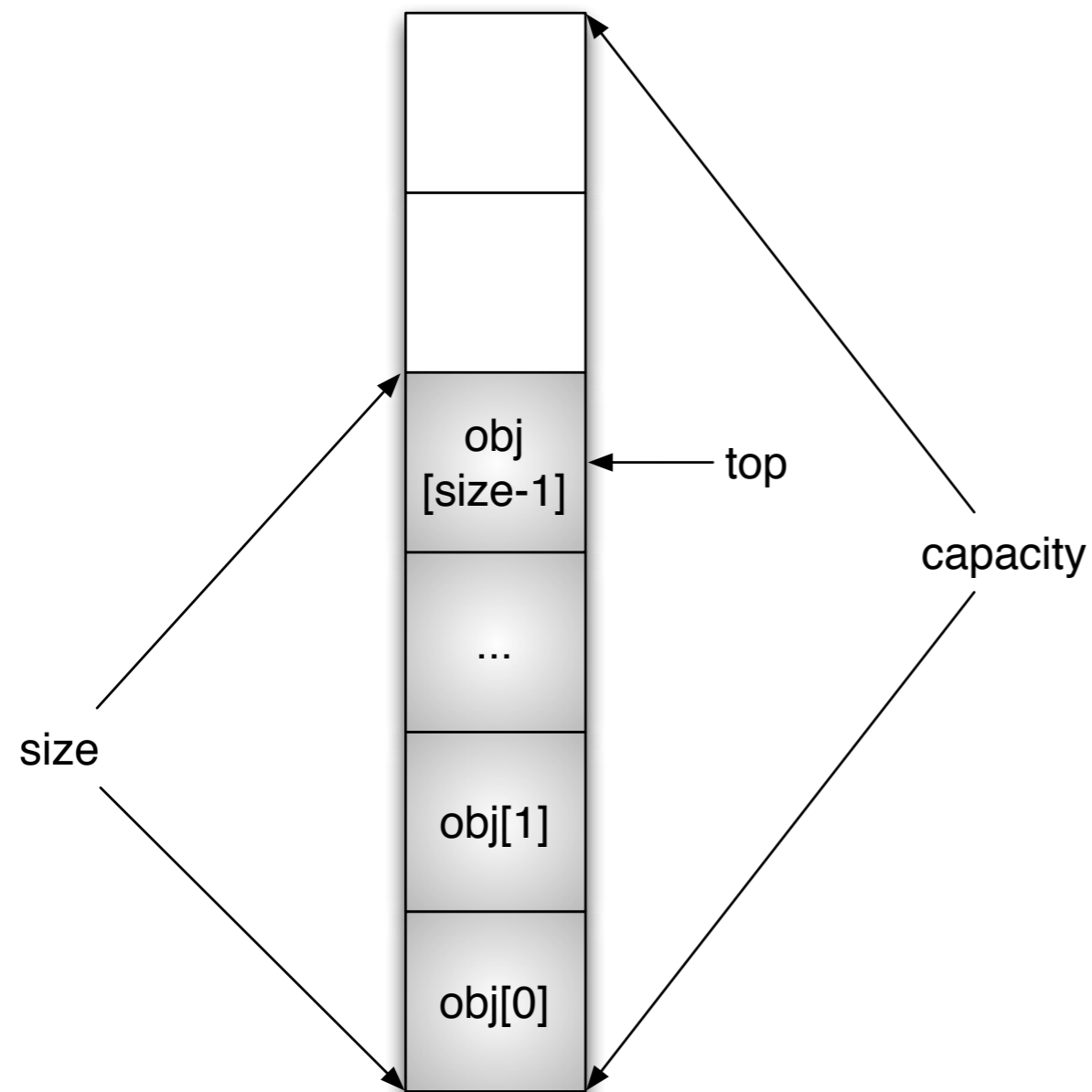
$$\text{pop}(\text{push}(x, s)) = s$$

Implementing a stack



- using an C-array `obj`
- array length defines the *capacity* of the stack
- *size* denotes the number of elements in stack
- `obj[size-1]` is the top element

Implementing a stack



```
struct Stack
{
    value_type* obj;
    size_type  size;
    size_type  capacity;
};

typedef struct Stack Stack;
```

Implementing `push_stack`

```
// Inserts v at the top of the stack.  
//  
void push_stack(Stack* stack, value_type v)  
{  
    if (!full_stack(stack))  
        stack->obj[stack->size++] = v;  
}
```

Implementing `pop_stack`

```
// Removes the element at the top of the stack.  
// Requires that the stack is not empty  
//  
void pop_stack(Stack* stack)  
{  
    if (!empty_stack(stack))  
        --stack->size;  
}
```


Verification of Stack

- implementing a stack is not difficult
- however, you are responsible to verify that your implementation is correct
- functional verification must be based on stack axioms

From Unit Tests to Unit Proofs

A Unit Test for Axiom 3

```
void stack_axiom3(Stack* s, Stack* t, value_type v)
{
    CU_ASSERT(equal_stack(s, t));

    push_stack(s, v);
    pop_stack(s);

    CU_ASSERT(equal_stack(s, t));
}
```

Tests need Test Objects

```
stack_axiom3(Stack* s, Stack* t, value_type v);
```

- must be executed with various values
 - e.g. empty, full, and “normal” stacks
- what is sufficient testing?

Unit Proofs

- proofs can handle very large, even infinite state spaces
- but requires formal specification

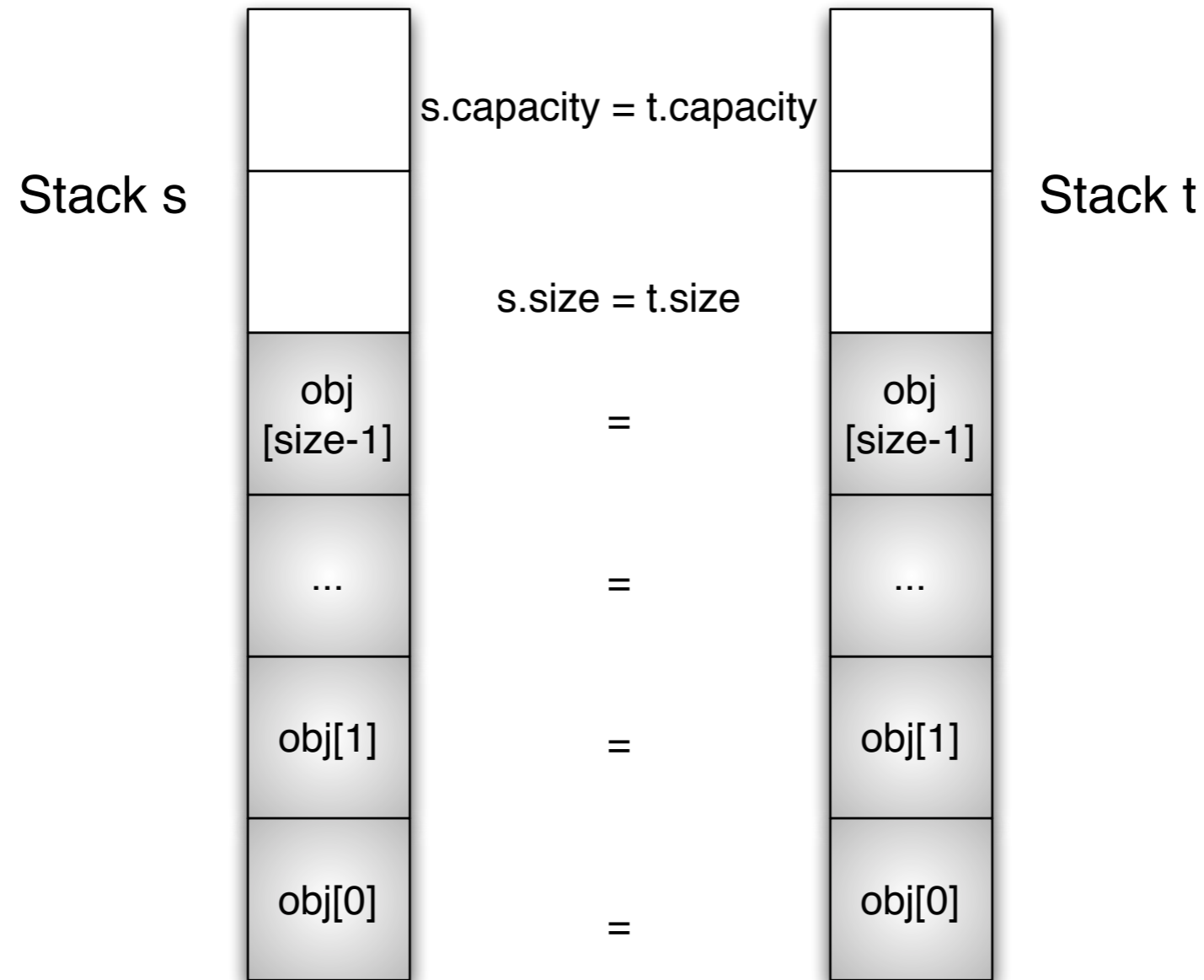
Axiom 3 revisited

```
void stack_axiom3(Stack* s, Stack* t, value_type v)
{
    CU_ASSERT(equal_stack(s, t));

    push_stack(s, v);
    pop_stack(s);

    CU_ASSERT(equal_stack(s, t));
}
```

Equality of stacks



Equality of stacks s and t

```
predicate IsEqualStack(Stack* s, Stack* t) =  
  
    s->capacity == t->capacity &&  
  
    s->size == t->size &&  
  
    \forall integer i;  
        0 <= i < s->size ==> s->obj[i] == t->obj[i];
```

IsEqualStack should be

- reflexive: $s == s$
- symmetric: $s == t \Rightarrow t == s$

- transitive:

$$(s == t \ \&\& \ t == u) \Rightarrow s == u$$

- an *equivalence relation*

Formalizing Reflexivity

```
lemma IsEqualStackReflexive :
```

```
  \forall Stack* s; IsEqualStack(s, s);
```

Formalizing Symmetry

```
lemma IsEqualStackSymmetric :
```

```
  \forall Stack *s1, *s2;
```

```
    IsEqualStack(s1, s2) ==> IsEqualStack(s2, s1);
```

Formalizing Transitivity

```
lemma IsEqualStackTransitive :
```

```
  \forall Stack *s1, *s2, *s3;
```

```
    IsEqualStack(s1, s2) && IsEqualStack(s2, s3)  
    ==> IsEqualStack(s1, s3);
```

Automatic Verification of these Lemmas

File	Configuration	Proof				
Proof obligations			Alt-Ergo 0.91	Simplify 1.5.7	Yices 1.0.28 (SS)	CVC3 2.2 (SS)
▼ User goals			✓	✓	✓	✓
Lemma IsEqualStackReflexive			●	●	●	●
Lemma IsEqualStackSymmetric			●	●	●	●
Lemma IsEqualStackTransitive			●	●	●	●

Using the equality predicate in Axiom 3

```
/*@  
    requires IsEqualStack(s, t);  
  
    assigns *s;  
  
    ensures IsEqualStack(s, t);  
*/  
void stack_axiom3(Stack* s, Stack* t, value_type v)  
{  
    push_stack(s, v);  
    pop_stack(s);  
}
```

Not true for full stacks

```
/*@  
    requires IsEqualStack(s, t);  
  
    assigns *s;  
  
    ensures IsEqualStack(s, t);  
*/  
void stack_axiom3(Stack* s, Stack* t, value_type v)  
{  
    push_stack(s, v);  
    pop_stack(s);  
}
```


A predicate for full stacks

predicate

```
IsFullStack{L}(Stack* s) =  
(s->size == s->capacity);
```

Excluding full stacks

```
/*@  
    requires IsEqualStack(s, t);  
    requires !IsFullStack(s);  
  
    assigns *s;  
  
    ensures IsEqualStack(s, t);  
*/  
void stack_axiom3(Stack* s, Stack* t, value_type v)  
{  
    push_stack(s, v);  
    pop_stack(s);  
}
```

What is a valid stack?

```
predicate IsValidStack{L}(Stack* s) =  
  
    \valid(s) &&  
  
    0 < s->capacity &&  
  
    0 <= s->size <= s->capacity &&  
  
    IsValidRange(s->obj, s->capacity);
```

Final Form of Axiom 3

```
/*@
    requires IsValidStack(s);
    requires IsValidStack(t);

    requires IsEqualStack(s, t);
    requires !IsFullStack(s);

    assigns *s;

    ensures IsEqualStack(s, t);
*/
void stack_axiom3(Stack* s, Stack* t, value_type v)
{
    push_stack(s, v);
    pop_stack(s);
}
```

Can we prove Axiom 3?

- no!
- we have not yet formally specified `push_stack` and `pop_stack`

```

/*@
  requires IsValidStack(s);
  ensures  IsValidStack(s);

  behavior not_empty:
    assumes !IsEmptyStack(s);
    assigns s->size;
    ensures s->size == \old(s->size) - 1;
    ensures !IsFullStack(s);

  behavior empty:
    assumes IsEmptyStack(s);
    assigns \nothing;

  complete behaviors;
  disjoint behaviors;
*/
void pop_stack(Stack* s);

```

Verification of pop_stack

Proof obligations	Alt-Ergo 0.92	Simplify 1.5.7	Z3 2.1 (SS)	Yices 1.0.28 (SS)	CVC3 2.2 (SS)	Statistics
▷ User goals	✓	✓		✓	✓	3/3
▷ Function pop_stack Default behavior	✓	✓		✗	✓	3/3
▷ Function pop_stack Normal behavior `empty`	✓	✓		✓	✓	2/2
▷ Function pop_stack Normal behavior `not_empty`	✓	✓		✓	✗	6/6
▷ Function pop_stack Safety	✓	✓		✓	✓	4/4

```

/*@
  requires IsValidStack(stack);
  ensures  IsValidStack(stack);

  behavior not_full:
    assumes !IsFullStack(stack);
    assigns stack->size;
    assigns stack->obj[stack->size];
    ensures stack->size == \old(stack->size) + 1;
    ensures stack->obj[\old(stack->size)] == v;
    ensures !IsEmptyStack(stack);

  behavior full:
    assumes IsFullStack(stack);
    assigns \nothing;

  complete behaviors;
  disjoint behaviors;
*/
void push_stack(Stack* stack, value_type v);

```


Verification of push_stack

Proof obligations	Alt-Ergo 0.92	Simplify 1.5.7	Z3 2.1 (SS)	Yices 1.0.28 (SS)	CVC3 2.2 (SS)	Statistics
▷ User goals	✓	✓		✓	✓	3/3
▷ Function push_stack Default behavior	✓	✓		✗	✓	3/3
▷ Function push_stack Normal behavior `full`	✓	✓		✓	✓	4/4
▷ Function push_stack Normal behavior `not_full`	✗	✓		✗	✗	10/10
▷ Function push_stack Safety	✓	✓		✗	✓	6/6

Verification of Axiom 3

```
/*@
    requires IsValidStack(s);
    requires IsValidStack(t);

    requires IsEqualStack(s, t);
    requires !IsFullStack(s);

    assigns *s;

    ensures IsEqualStack(s, t);
*/
void stack_axiom3(Stack* s, Stack* t, value_type v)
{
    push_stack(s, v);
    pop_stack(s);
}
```

Verification of Axiom 3

Proof obligations	Alt-Ergo 0.91	Simplify 1.5.7	Yices 1.0.28 (SS)	CVC3 2.2 (SS)
▷ User goals	✓	✓	✓	✓
▷ Function stack_axiom3 Default behavior	✗	✓	✗	✗
▷ Function stack_axiom3 Safety	✓	✓	✗	✓

Verification of Axiom 3

We have proved that axiom 3 is satisfied by all valid stacks that are not full

Compare this with unit tests!

Representation Independence (RI)

RI is well known to Java Programmers

- consider `java.util.Stack`
- overrides `equal`, `hashCode`, and `toString`
- if `s.equal(t)` then
 `s.hashCode() == t.hashCode()`
must hold
- RI-violations are a common source of errors in Java

Leibniz Law

Two or more objects are identical if they have all their properties in common.

RI particularly means

```
//@ assert IsEqualStack(s, t);
```

```
    push_stack(s, v);  
    push_stack(t, v);
```

```
//@ assert IsEqualStack(s, t);
```

More precisely

```
/*@
    requires IsValidStack(s);
    requires IsValidStack(t);
    requires IsEqualStack(s, t);

    assigns *s;
    assigns *t;

    ensures IsEqualStack(s, t);
*/
void push_stack_independence(Stack* s, Stack* t, value_type v)
{
    push_stack(s, v);
    push_stack(t, v);
}
```

Another Example

```
/*@  
    requires IsValidStack(s);  
    requires IsValidStack(t);  
    requires IsEqualStack(s, t);  
  
    assigns \nothing;  
  
    ensures IsEqualStack(s, t);  
    ensures \result == 1;  
*/  
bool top_stack_independence(Stack* s, Stack* t)  
{  
    return top_stack(s) == top_stack(t);  
}
```

Verification of RI

Proof obligations	Alt-Ergo 0.91	Simplify 1.5.7	Yices 1.0.28 (SS)	CVC3 2.2 (SS)
▷ User goals	✓	✓	✓	✓
▷ Function empty_stack_independe Default behavior	✓	✓	✗	✓
▷ Function empty_stack_independe Safety	✓	✓	✗	✓
▷ Function full_stack_independen Default behavior	✓	✓	✗	✓
▷ Function full_stack_independen Safety	✓	✓	✗	✓
▷ Function pop_stack_independenc Default behavior	✗	✓	✗	✓
▷ Function pop_stack_independenc Safety	✓	✓	✗	✓
▷ Function push_stack_independen Default behavior	✗	✓	✗	✗
▷ Function push_stack_independen Safety	✓	✓	✗	✓
▷ Function top_stack_independenc Default behavior	✓	✓	✗	✓
▷ Function top_stack_independenc Safety	✓	✗	✗	✓

Conclusions

- not just replacing *unit tests* by *unit proofs*
- a much closer look on equality and representation independence

Questions please!

Implementing `empty_stack`

```
// Returns true if the stack contains no elements,  
// and false otherwise  
//  
bool empty_stack(const Stack* stack)  
{  
    if (stack->size == 0)  
        return 1;  
    else  
        return 0;  
}
```


Implementing `full_stack`

```
// Returns true if the stack contains
// as many elements as its capacity
// and false otherwise
//
bool full_stack(const Stack* stack)
{
    return (stack->size == stack->capacity) ? 1 : 0;
}
```

Implementing `top_stack`

```
// Returns the element at the top of the stack
//
value_type top_stack(const Stack* stack)
{
    if (!empty_stack(stack))
        return stack->obj[stack->size - 1];
    else
        return 0;
}
```