

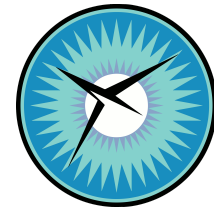
# Compact Position Reporting Algorithm

A verified floating-point implementation in C

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# The Algorithm

# The ADS-B System



- Automatic Dependent Surveillance - Broadcast
  - Supports NextGen
    - Next generation of air traffic management systems
  - Aircraft periodically *broadcasts* accurate surveillance information to ground stations and near aircraft
    - position and velocity
  - *Automatic* – no pilot intervention needed
  - *Dependent* – on navigation system
- Mandatory on Jan 1, 2020 (in USA and Europe)
  - More than 40000 aircraft currently equipped

# The ADS-B Protocol



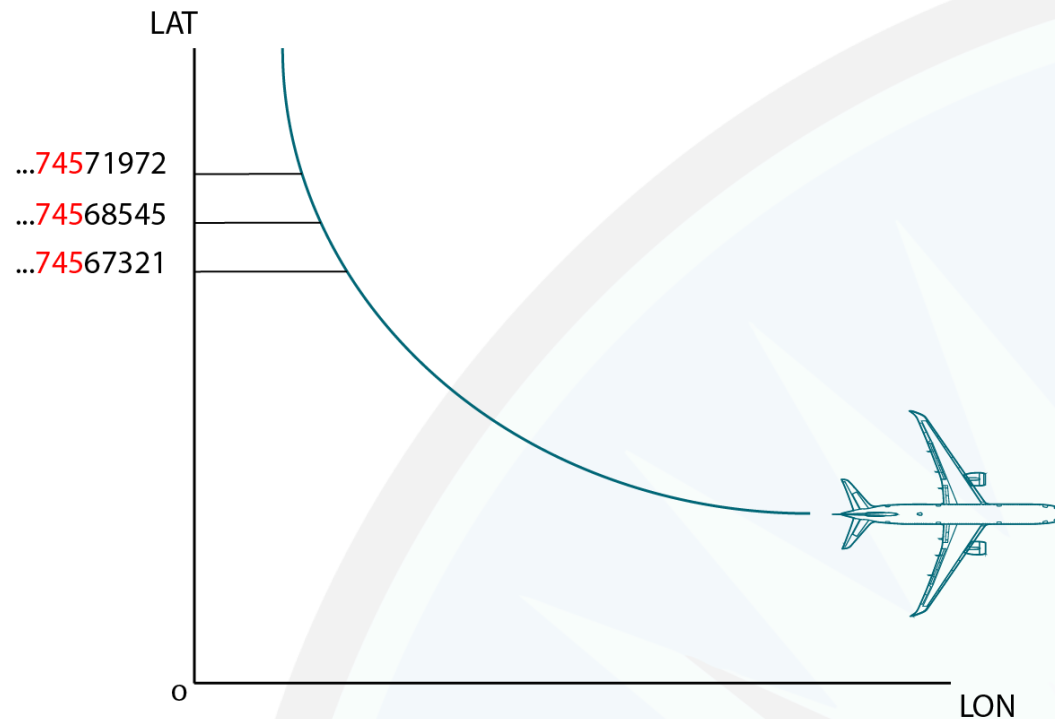
- Pros: broadcast vs. radar-based approaches
  - ✓ More precise
    - NextGen requirement: position granularity of ~5.1 meters
  - ✓ More coverage
- Cons: Make use of existent hardware
  - ✗ TCAS transponders
  - ✗ 35 bits for position data in the broadcast message
  - ✗ Too coarse granularity (~300 meters)
    - if raw positions are transmitted



# Compact Position Reporting



Contiguous transmitted positions share prefixes



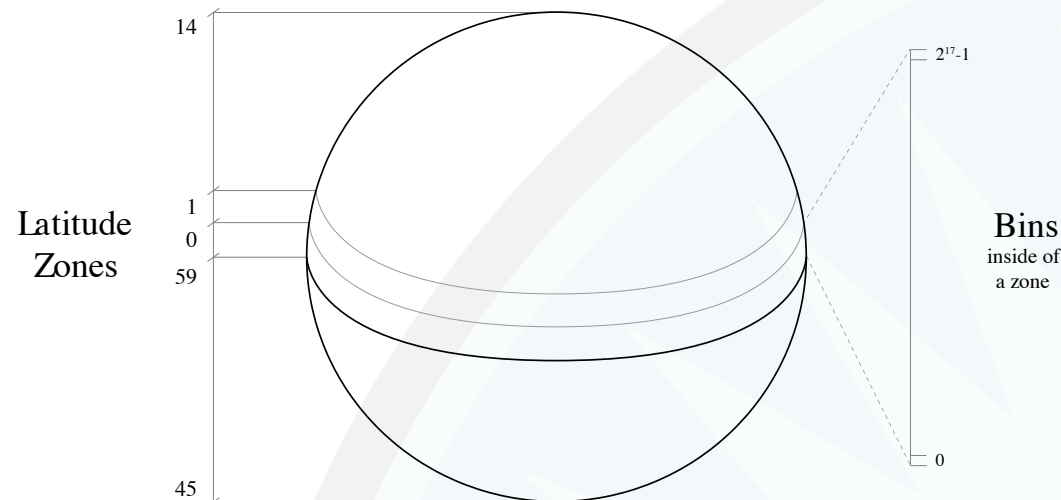
Idea: transmit only 17 less significant bits

# Focus on Latitude First

# Latitude Zones

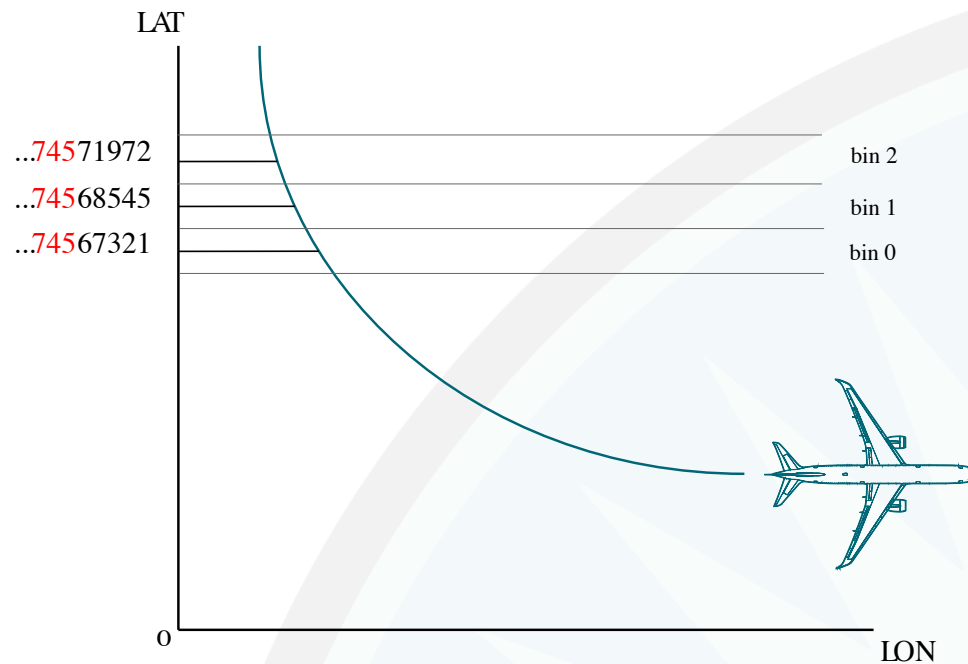


- Divide the globe into 60 equally sized zones
- Divide each zone in  $2^{17}$  bins



Zone Size:  $D_{lat} = 360/60 = 6$  degrees

# Reported Latitude



Broadcast only the corresponding *bin number* (YZ)

# Encoding Latitude



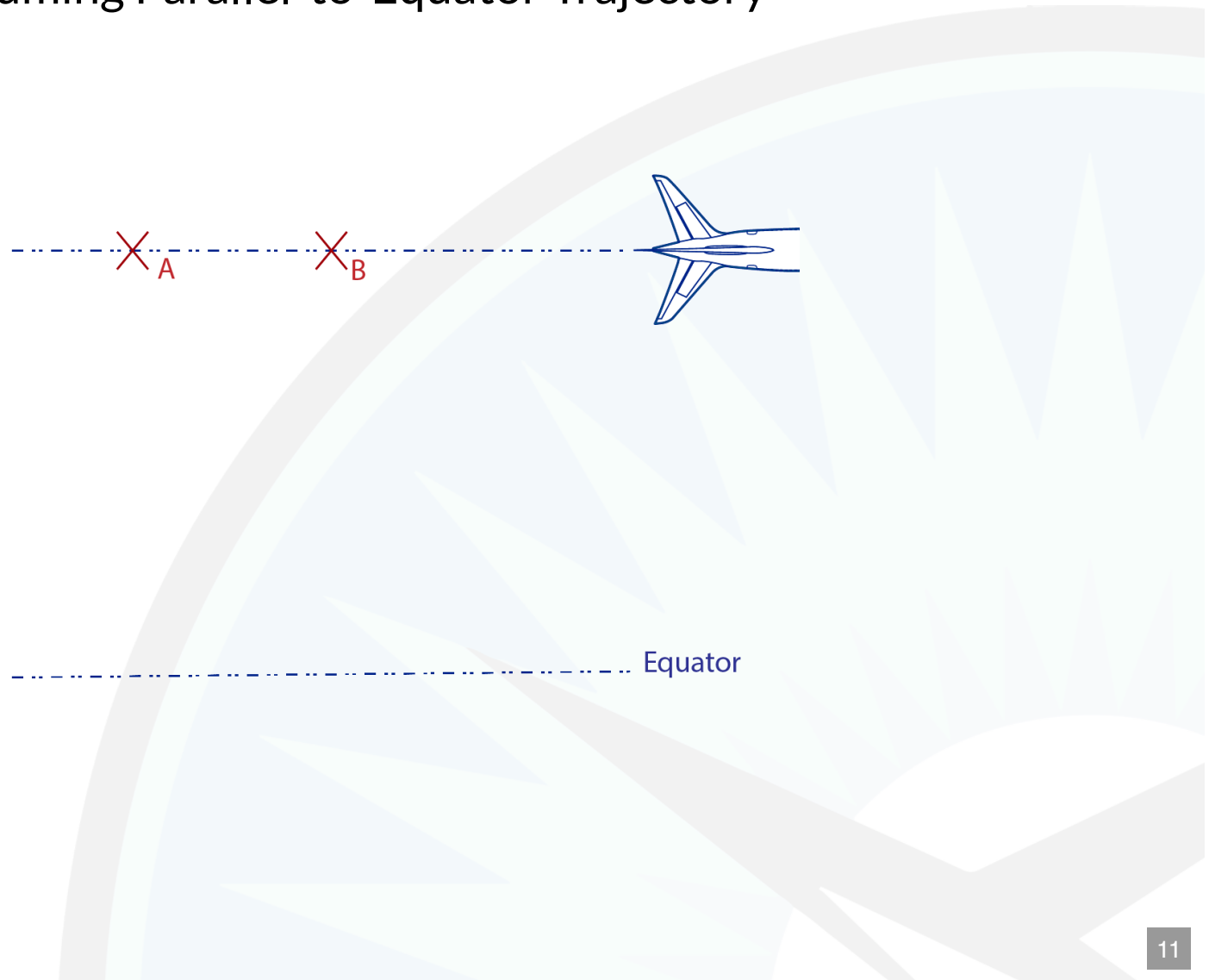
- To encode  $lat$ , calculate:
  1. Distance from southern edge of enclosing zone
    - $\text{mod}(lat, Dlat)$
  2. Proportion w.r.t. the entire zone
    - $\text{mod}(lat, Dlat) \cdot \frac{1}{Dlat}$
  3. Correspondent *bin* number
    - $\text{mod}(lat, Dlat) \cdot \frac{1}{Dlat} \cdot 2^{17}$
  4. Round to the nearest integer
    - $ZY = \left\lfloor \text{mod}(lat, Dlat) \cdot \frac{1}{Dlat} \cdot 2^{17} + \frac{1}{2} \right\rfloor$

# How to Recover the Zone Index

# Recovering Zone Index



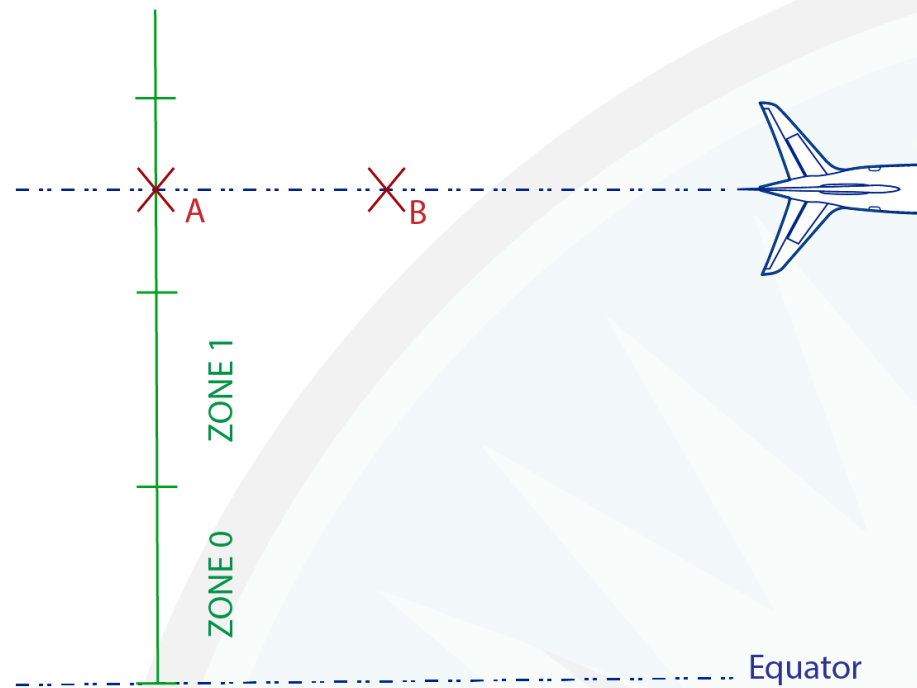
Assuming Parallel-to-Equator Trajectory



# Recovering Zone Index



Assuming Parallel-to-Equator Trajectory

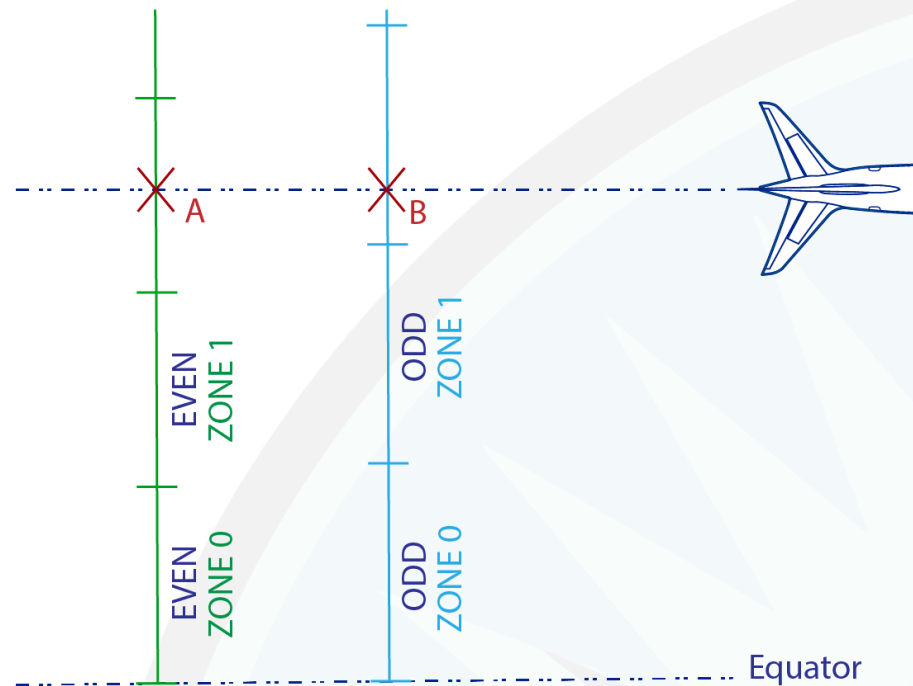




# Recovering Zone Index



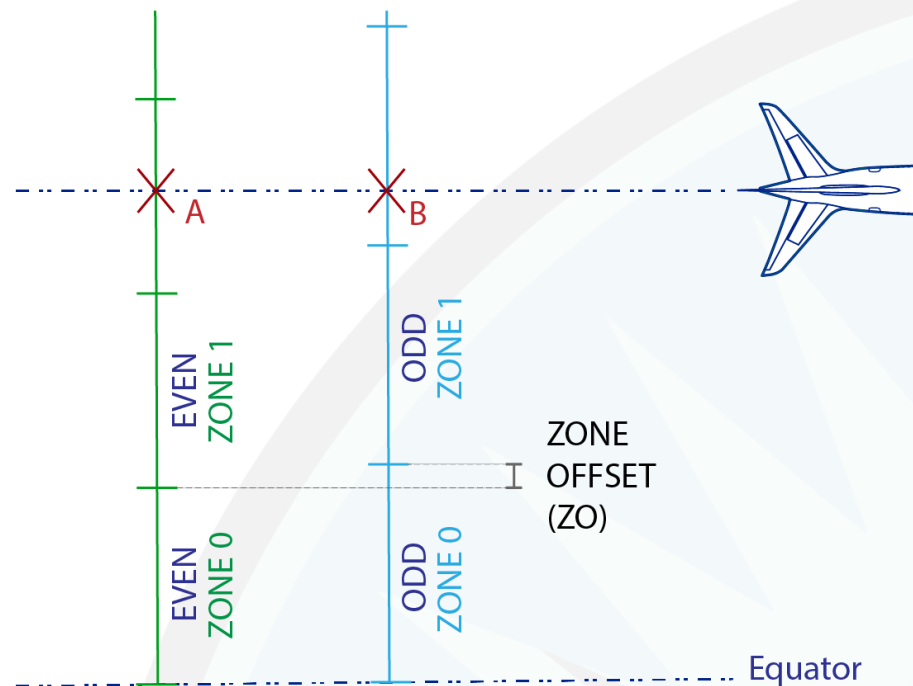
Assuming Parallel-to-Equator Trajectory



# Recovering Zone Index



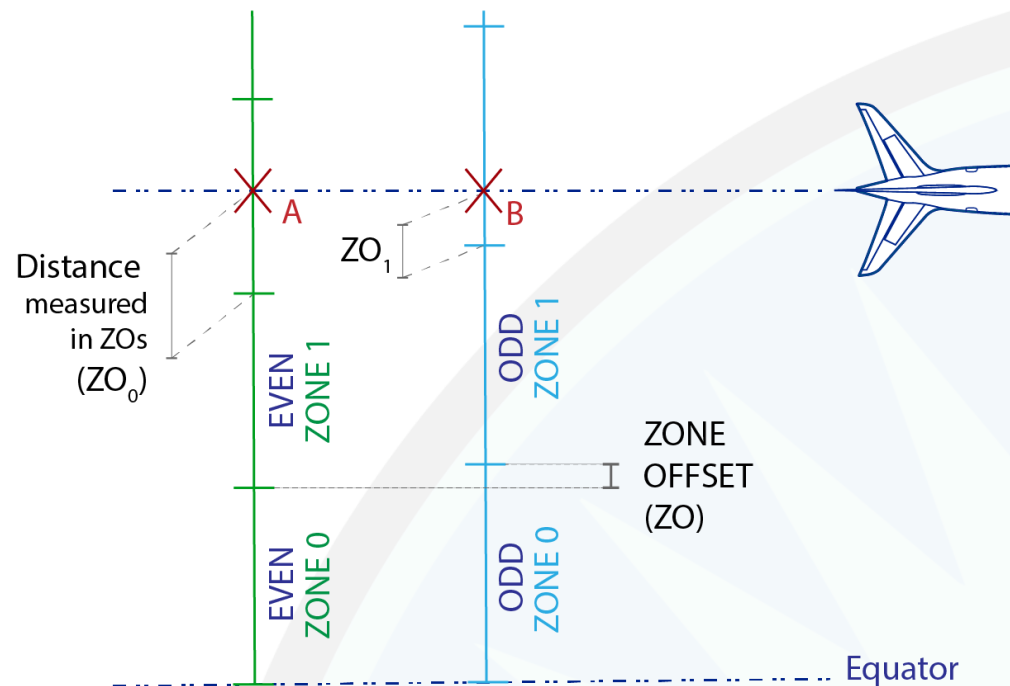
Assuming Parallel-to-Equator Trajectory



# Recovering Zone Index



Assuming Parallel-to-Equator Trajectory



$$\text{Zone Index: } ZI := \lfloor ZO_0 - ZO_1 + 1/2 \rfloor$$



## Relaxing parallel-to-the-Equator restriction

- According to the standard, if two latitudes  $A$  and  $B$  are less than half zone offset apart from each other,
  - $A$  and  $B$  lie in the same zone, or
  - $A$  is one zone ahead w.r.t.  $B$
- To deal with both cases

$$ZI = \begin{cases} \text{mod}(\lfloor ZO_0 - ZO_1 + 1/2 \rfloor, 60) & \text{even zone index} \\ \text{mod}(\lfloor ZO_0 - ZO_1 + 1/2 \rfloor, 59) & \text{odd zone index} \end{cases}$$



Given an *even* and an *odd* bin number  $YZ_0$  and  $YZ_1$ , the recovered latitude  $Rlat_i$  is defined as

$$Rlat_i(YZ_0, YZ_1) := Dlat_i \left( \text{mod} (\lfloor ZO_0 - ZO_1 + 1/2 \rfloor, 60 - i) + YZ_i \frac{1}{2^{17}} \right)$$

where  $ZO_i$  (zone offset)  $ZO_i := \frac{Dlat_i}{ZO} \cdot \frac{YZ_i}{2^{17}}$  where  $i \in \{0, 1\}$

- Note that
  - $Rlat_i$  returns the center of the *bin* where the input latitude lies.
  - Decoded latitude is at most at half-bin size from the input latitude

# What About Longitudes?

# Dealing with Longitudes



- Goal: same encoding resolution everywhere
  - as close to a constant as possible all around the globe
- Same idea
  - ~Equally sized zones divided in  $2^{17}$  bins
- One distinctive feature
  - Longitude (radial) size shrinks when approaching the poles
  - Number of longitude zones is a function of latitude
    - reducing the number of zones as latitude increases

# NL Function



- $NL(lat)$ : number of even longitude zones at latitude  $lat$

$$NL(lat) = \begin{cases} 59 & \text{if } lat = 0, \\ \left\lfloor 2\pi \left( \arccos \left( 1 - \frac{1 - \cos\left(\frac{\pi}{30}\right)}{\cos^2\left(\frac{\pi}{180} |lat|\right)} \right) \right)^{-1} \right\rfloor & \text{if } |lat| < 87, \\ 2 & \text{if } |lat| = 87, \\ 1 & \text{if } |lat| > 87. \end{cases}$$

- In practice, computing this function is inefficient
  - A lookup table of transition latitudes is pre-calculated



# Global Decoding



- Latitude, given two encoded latitudes

$$\text{Rlat}_i(\text{YZ}_0, \text{YZ}_1) := \text{Dlat}_i \left( \text{mod} \left( \left\lfloor \frac{59\text{YZ}_0 - 60\text{YZ}_1}{2^{17}} + \frac{1}{2} \right\rfloor, 60 \right) + \frac{\text{YZ}_i}{2^{17}} \right)$$

- Longitude, given two encoded positions

$$\text{Rlon}_i(\text{YZ}_0, \text{YZ}_1, \text{XZ}_0, \text{XZ}_1) := \text{Dlon}_i \left( \text{mod} \left( \left\lfloor \frac{(nl-1)\text{XZ}_0 - nl \cdot \text{XZ}_1}{2^{17}} + \frac{1}{2} \right\rfloor, nl'_i \right) + \frac{\text{XZ}_i}{2^{17}} \right)$$

where

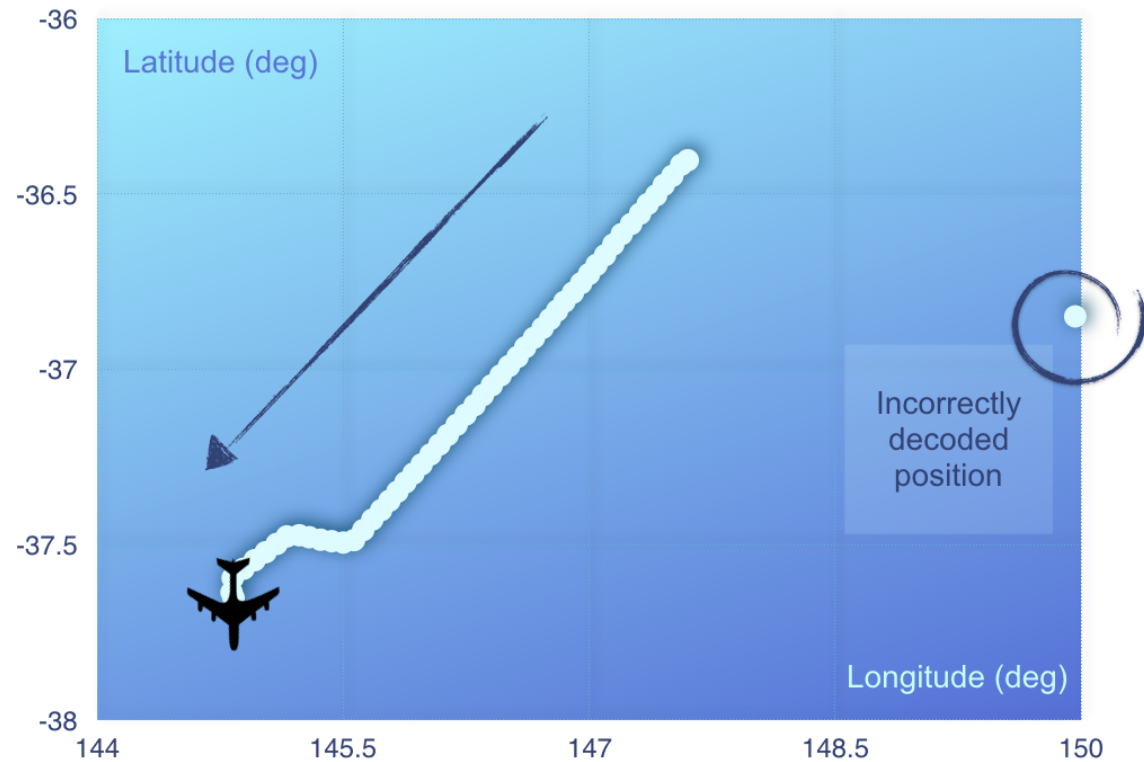
- $nl := \text{NL}(\text{Rlat}_0(\text{YZ}_0, \text{YZ}_1))$ , must be equal to  $\text{NL}(\text{Rlat}_1(\text{YZ}_0, \text{YZ}_1))$
- $nl'_i := \max(nl - i, 1)$ , since  $nl$  is 1 if  $|\text{Rlat}_i(\text{YZ}_0, \text{YZ}_1)| > 87$
- $\text{Dlon}_i := 360/nl'_i$

# Local Decoding



- Additional decoding method
- Uses a *reference position* and one position message
  - instead of two position messages
- Positions apart for no more than half of a zone
  - According to the standard
  - Allows for bigger separation between received positions
- Idea: create a sliding region 1 zone wide
  - Centered on reference position
  - Each bin number occurs only once in the region

# Known Issues



Reported by Airservices Australia (2007)

# Analysis of the Algorithm



- Accomplishments:

1. Found technical issues in the standard

- Counterexamples for the real-valued model

2. Amended version proven correct

- Prototype Verification System (PVS)

3. Proposed simpler formulation

- reducing numerical complexity

4. Prototype implementation formally verified

- C, PVS, Frama-C, Gappa, Alt-Ergo

**Dutle A., Moscato M., Titolo L., Muñoz C.** *A Formal Analysis of the Compact Position Reporting Algorithm.* VSTTE 2017.

**Titolo L., Moscato M., Muñoz C., Dutle A., Bobot F.** *A Formally Verified Floating-Point Implementation of the Compact Position Reporting Algorithm.* FM 2018.

# Technical Issues



- Counterexamples found for both decoding settings
  - Even Assuming (exact) real-valued arithmetics
  - For example, in the *global decoding* case
    - $\text{lat}_0 = 363373617 \cdot 360/2^{32} \approx 30.4576247279$
    - $\text{lat}_1 = 363980245 \cdot 360/2^{32} \approx 30.5084716994$
    - decoded positions are further away for more than a *bin*
- Correctness proved on tightened requirements
  - max. distance of input positions decreased by half-bin size

# Numerical Simplifications



- Mathematically equivalent expressions suggested
  - Numerically simpler
  - Equivalence formally proven
- Example: equivalent calculation of NL lookup table
  - removing four operations in total
  - $\text{lat}_{\text{NL}}(nl) := \frac{180}{\pi} \arccos\left(\frac{\sin(\pi/60)}{\sin(\pi/nl)}\right)$ .
- Example: cancellation instead of division
  - Reducing complexity of encoding algorithm
  - $\frac{\text{mod}(a,b)}{b} = \frac{a - b * \lfloor \frac{a}{b} \rfloor}{b} = \frac{a}{b} - \lfloor \frac{a}{b} \rfloor$

# Example: Latitude Global Decoding



- According to the standard:

$$\text{Rlat}_0(\text{YZ}_0, \text{YZ}_1) := \text{Dlat}_0 \left( \text{mod} \left( \left\lfloor \frac{59\text{YZ}_0 - 60\text{YZ}_1}{2^{17}} + \frac{1}{2} \right\rfloor, 60 \right) + \frac{\text{YZ}_0}{2^{17}} \right)$$

- Simplified version of global decoding (i=0) in ACSL

```
/*@ axiomatic real_function {  
  logic real rLatr (int yz0,int yz1) =  
    \let dLatr = 360.0 / 60.0;  
    \let jar   = (59.0*yz0 - 60.0*yz1 + 0x1.0p+16)*0x1.0p-17;  
    \let jr    = \floor(jar);  
    \let j60ir = jr/60.0;  
    dLatr*((jr-60.0*(\floor(j60ir)))+yz0*0x1.0p-17); } @*/
```

# Example: Latitude Global Decoding



- Simplified version of global decoding (i=0) in ACSL

```
/*@ axiomatic real_function {  
  logic real rLatr (int yz0,int yz1) =  
    \let dLatr = 360.0 / 60.0;  
    \let jar   = (59.0*yz0 - 60.0*yz1 + 0x1.0p+16)*0x1.0p-17;  
    \let jr    = \floor(jar);  
    \let j60ir = jr/60.0;  
    dLatr*((jr-60.0*\floor(j60ir))+yz0*0x1.0p-17); } @*/
```

- Translated by hand into a PVS declaration

```
rLatr_i_0 (yz0,yz1:int): real =  
  LET dLatr = 360 / 60 IN  
  LET jar   = (59*yz0 - 60*yz1 + 2^16) * 2^-17 IN  
  LET jr    = floor(jar) IN  
  LET j60ir = jr/60 IN  
  dLatr * ((jr - 60*floor(j60ir)) + yz0 * 2^-17)
```

- Proven to be equivalent to version from the standard



# Example: Latitude Global Decoding



```
/*@ requires 0 <= yz0 <= 131071; requires 0 <= yz1 <= 131071;
   requires \floor(yz0) == yz0; requires \floor(yz1) == yz1;
   ensures \abs(\result - rLatr(yz0,yz1)) <= 0.000022888; */
fp rLatf (int yz0, int yz1) {
  fp res, rLat1; fp dLatf = 360.0 / 60.0;
  fp j1f = (59.0 * yz0 - 60.0 * yz1 + 0x1.0p+16) * 0x1.0p-17;
  /*@ assert j1f:
     \let j1r = (59.0 * yz0 - 60.0 * yz1 + 0x1.0p+16) *0x1.0p-17;
     j1f == j1r; */
  fp jf = floor(j1f);
  /*@ assert jf:
     \let j1r = (59.0 * yz0 - 60.0 * yz1 + 0x1.0p+16) *0x1.0p-17;
     \let jr = \floor(j1r);
     jf == jr; */
  /*@ assert values_for_jf: -60.0 <= jf <= 59.0; */
  /*@ assert jf represents an integer: \floor(jf) == jf; */
}
```

Frama-C/WP & Alt-Ergo+Gappa: the floating-point result is at most  $0.000022888^\circ$  (half bin size) apart from the logical result.

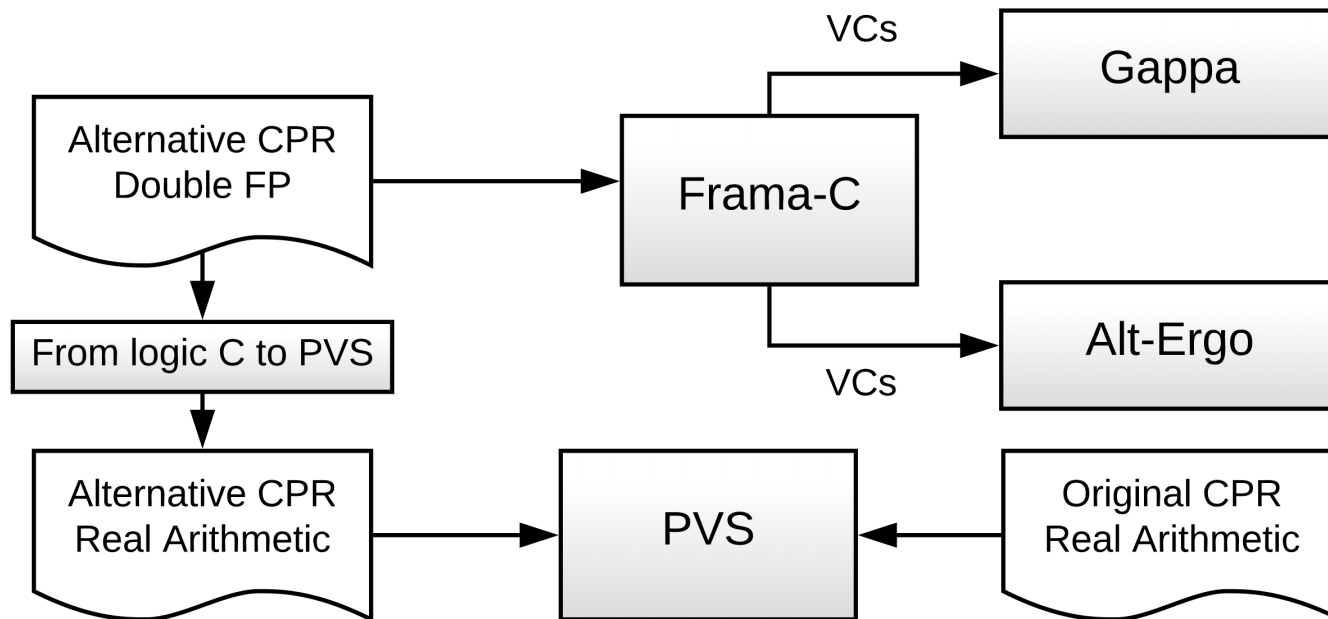
## Result of the Verification Process



Floating-point version has the expected granularity:  
decoded and input positions are less than  $\frac{1}{2}$  bin apart

- Amended CPR version has been proved correct, i.e.,
  - decoded latitude lies in the center of a bin and
  - it is less than half bin apart from the input
- It coincides with the ACSL logic definition
- C version is less than half bin size apart from it

# Verification Approach



- logic ACSL declarations translated to PVS by hand
- proved equivalent to existent CPR formalization
- C code verified using Frama-C/WP/Alt-Ergo/Gappa

# Concluding Remarks



- Synergetic use of diverse analysis tools on
  - complex verification effort
  - relatively simple algorithm
    - no loops, no pointers, no arrays
- Proposed algorithm is being considered as reference implementation of CPR
  - RTCA DO-260B/Eurocae ED-102A

# Future Work



- Extend results to other CPR modalities
  - Airborne, Surface, Coarse TIS-B
- Develop CPR integer-valued version
  - correctness (PVS) + verified implementation (Frama-C)
- Analysis of Floating-Point Programs
  - Frama-C: WP plugin to export VCs directly to PVS
  - Floating-point programs: Frama-C + PRECiSA
    - <http://precisa.nianet.org/>

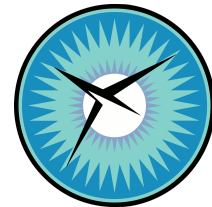
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Thank you for you attention